

# Information Theory

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## Today's Topics

- Entropy
- Conditional Entropy
- Example
- Joint Entropy
- Chain Rule

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## Entropy H(S)

- Entropy is the average information content of a source

$$H(S) = E[I(s_k)]$$
$$H(S) = \sum_{k=0}^{K-1} p_k \log_2 \left( \frac{1}{p_k} \right)$$

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## Entropy

### Comments

Entropy is a measure of how much information is encoded in a message. Higher the entropy, higher the information content.

We could also say entropy is a measure of uncertainty in a message.

Information and uncertainty are equivalent concepts.

Entropy gives the actual number of bits of information contained in a message source.

Example: if the probability of the character `e` appearing in this slide is 1/16, then the information content of this character is 4 bits.

So the character string `eeee` has a total of 20 bits (contrast this to using an 8-bit ASCII coding that could result in 40 bits to represent `eeee`).

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## Conditional Entropy

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## Conditional Entropy

$$H(Y|X=v)$$

Suppose I'm trying to predict output Y and I have input X

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### Conditional Entropy $H(Y|X=v)$

**Example** If I know the student's major could I predict if he likes computer games?  
**Input:** X = College Major  
**Output:** Y = Likes "Computer Games"

X	Y
Math	Yes
History	No
CS	Yes
Math	No
Math	No
CS	Yes
History	No
Math	Yes

X	Y	Yes	No	
Math	Yes	0.25	0.25	$P(X=Math)=0.5$
CS	Yes	0.25	0	$P(X=CS)=0.25$
History	No	0	0.25	$P(X=History)=0.25$
		$P(Y=Yes)=0.5$	$P(Y=No)=0.5$	1

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### Conditional Entropy $H(Y|X=v)$

**Example** If I know the student's major could I predict if he likes computer games?  
**Input:** X = College Major  
**Output:** Y = Likes "Computer Games"

X	Y
Math	Yes
History	No
CS	Yes
Math	No
Math	No
CS	Yes
History	No
Math	Yes

X	Y	Yes	No	
Math	Yes	0.25	0.25	$P(X=Math)=0.5$
CS	Yes	0.25	0	$P(X=CS)=0.25$
History	No	0	0.25	$P(X=History)=0.25$
		$P(Y=Yes)=0.5$	$P(Y=No)=0.5$	1

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### Specific Conditional Entropy $H(Y|X=v)$

**Example** **Definition of Specific Conditional Entropy:**  
 $H(Y|X=v) =$  The entropy of Y among only those records in which X has value v

X	Y
Math	Yes
History	No
CS	Yes
Math	No
Math	No
CS	Yes
History	No
Math	Yes

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### Specific Conditional Entropy $H(Y|X=v)$

**Example** **Definition of Specific Conditional Entropy:**  
 $H(Y|X=v) =$  The entropy of Y among only those records in which X has value v

**Example:**  
 $H(Y|X=Math) = -\sum_{y \in \{Yes, No\}} p(y|X=Math) \log p(y|X=Math)$   
 $= -p(Yes|X=Math) \log p(Yes|X=Math) - p(No|X=Math) \log p(No|X=Math)$   
 $= -0.5 \log 0.5 - 0.5 \log 0.5 = 1$

**$H(Y|X=Math) = 1$**

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### Specific Conditional Entropy $H(Y|X=v)$

**Example** **Definition of Specific Conditional Entropy:**  
 $H(Y|X=v) =$  The entropy of Y among only those records in which X has value v

**Example:**  
 $H(Y|X=History) = -\sum_{y \in \{Yes, No\}} p(y|X=History) \log p(y|X=History)$   
 $= -p(Yes|X=History) \log p(Yes|X=History) - p(No|X=History) \log p(No|X=History)$   
 $= -0 \log 0 - 1 \log 1 = 0$

**$H(Y|X=History) = 0$**

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### Specific Conditional Entropy $H(Y|X=v)$

**Example** **Definition of Specific Conditional Entropy:**  
 $H(Y|X=v) =$  The entropy of Y among only those records in which X has value v

**Example:**  
 $H(Y|X=CS) = -\sum_{y \in \{Yes, No\}} p(y|X=CS) \log p(y|X=CS)$   
 $= -p(Yes|X=CS) \log p(Yes|X=CS) - p(No|X=CS) \log p(No|X=CS)$   
 $= -1 \log 1 - 0 \log 0 = 0$

**$H(Y|X=CS) = 0$**

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## Specific Conditional Entropy $H(Y|X=v)$

### Example

X	Y
Math	Yes
History	No
CS	Yes
Math	No
Math	No
CS	Yes
History	No
Math	Yes

### Definition of Specific Conditional Entropy:

$H(Y|X=v)$  = The entropy of Y among only those records in which X has value v

### Example:

- $H(Y|X=Math) = 1$
- $H(Y|X=History) = 0$
- $H(Y|X=CS) = 0$

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## Conditional Entropy $H(Y|X)$

Is the amount of information contained in Y such that X is given

### Example

X	Y
Math	Yes
History	No
CS	Yes
Math	No
Math	No
CS	Yes
History	No
Math	Yes

### Definition of Conditional Entropy:

$H(Y|X)$  = The average specific conditional entropy of Y

= if you choose a record at random what will be the conditional entropy of Y, conditioned on that row's value of X

= Expected number of bits to transmit Y if both sides will know the value of X

$$= \sum_j \text{Prob}(X=v_j) H(Y|X=v_j) \quad 14$$

## Conditional Entropy

### Example

X	Y
Math	Yes
History	No
CS	Yes
Math	No
Math	No
CS	Yes
History	No
Math	Yes

### Definition of Conditional Entropy:

$H(Y|X)$  = The average conditional entropy of Y

$$= \sum_j \text{Prob}(X=v_j) H(Y|X=v_j)$$

### Example:

$v_j$	Prob( $X=v_j$ )	$H(Y X=v_j)$
Math	0.5	1
History	0.25	0
CS	0.25	0

$$H(Y|X) = 0.5 * 1 + 0.25 * 0 + 0.25 * 0 = 0.5$$

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## Joint Entropy

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## Joint Entropy

Is the amount of information contained in both events X and Y

$$H(X, Y) = -\sum_{x,y} p(x,y) \log p(x,y)$$

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## Joint Entropy

### Chain Rule

Relationship between conditional and joint entropy

$$H(X, Y) = H(X) + H(Y|X)$$

### Proof:

$$\begin{aligned} H(X, Y) &= -\sum_{x,y} P(x, y) \log_2 P(x, y) = -\sum_{x,y} P(x, y) \log_2 P(x) P(y|x) \\ &= -\sum_x P(x) \sum_y P(y|x) \log_2 P(x) - \sum_{x,y} P(x, y) \log_2 P(y|x) = H(X) + H(Y|X) \end{aligned}$$

$$H(X|Y) = -\sum_{x,y} P(x, y) \log_2 P(x|y)$$

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## Joint Entropy

Is the amount of information contained in both events X and Y

$$H(X,Y) = H(X) + H(Y|X)$$

Also  $H(X,Y) = H(Y) + H(X|Y)$

- Intuition: first describe Y and then X given Y
- From this:  $H(X) - H(X|Y) = H(Y) - H(Y|X)$

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## Joint Entropy

$$H(X,Y) = H(X) + H(Y|X)$$

Example

X	Y
Math	Yes
History	No
CS	Yes
Math	No
Math	No
CS	Yes
History	No
Math	Yes

$v_j$	Prob( $X=v_j$ )	$H(Y   X = v_j)$
Math	0.5	1
History	0.25	0
CS	0.25	0

$$H(X) = 1.5$$

$$H(Y) = 1$$

$$H(Y|X) = 0.5$$

$$H(X,Y) = H(X) + H(Y|X) = 1.5 + 0.5 = 2$$

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## Joint Entropy

Comments

$$\begin{aligned} H(X,Y) &= H(X) + H(Y|X) \\ &= H(Y) + H(X|Y) \end{aligned}$$

$$\begin{aligned} H(X,Y,Z) &= H(X) + H(Y|X) + H(Z|XY) \\ &\neq H(X) + H(Y) + H(Z) \end{aligned}$$

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