Brief Introduction to Prolog

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- Prolog data structures
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Tutorial 3



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Section 1

Introduction

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Introduction to Prolog Prolog Syntax Tutorial 1

Programming paradigms

Imperative translation from machine language to user commands Object object-oriented, versatile and recent Declarative detachment between program's goal and execution details Functional concerned with recursion, pattern matching, ... Logic focus on automatically reasoning about knowledge in the program

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Introduction to Prolog

 Prolog represents Horn Clauses - which are a subset of First Order Logics - where each clause can have at most one positive literal in the head.

 $\mathsf{head}{:}-\mathsf{body}{.}$

- A Prolog program is a set of **facts** and/or **rules** defining relations between entities.
- Facts represent relations which are assumed to be true (axioms)
- **Rules** can be true or false, depending on other relations in the program
- A Prolog program is resolved by computing the consequents of rules.

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Prolog syntax 1

Characters Letters and numbers Terms can be variables Person, _Father constants mary, 'The Family' compound family('mary&tom', date(25, feb, 1954)) Symbols :- ; , . Comments can be used for

line %Comment
text /* Also a comment */

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Prolog syntax 2

Operators + - * /

Unification = finds a set of variable substitutions that make two terms exacty the same.

Arithmetic Prolog does not evaluate mathematical expressions unless the operator **is** is used

> numbers is > < = =terms == =

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Family - facts

%Facts T1 mother_of(mary,john). mother_of(mary,anne). father_of(tom,anne). father_of(tom,john).

%Query T1 ?-mother_of(mary, Person).



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Observations

- There are two predicates in this example: **mother_of/2** and **father_of/2**.
- pred/N represents predicate pred(arg1, arg2, ..., argN), with *arity* of N arguments.
- These predicates have no body and so they are called facts.
- Facts can be seen as a multi-relational database in Prolog.

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Section 2

Prolog data structures

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Data structures: List

A few examples:

- []: empty list.
- [the,men, [like,to,fish]]
- [a,V1,b,[X,Y]]

List is composed of: **[Head**|**Tail] Head**: first element of the list (can be of any type). **Tail**: remaining elements as a list.

List	Head	Tail
[a,b,c]	a	[b,c]
[a]	a	[]
[]	no head	no tail

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Equality of Lists

- Unifications can happen inside lists
- The empty list does not have head or tail
- This is used as a stop criterion in list recursion

Examples:

[X,Y,Z]	=	[john,likes,fish]	X=john, Y=likes, Z=fish
[cat]	=	[X Y]	X=cat, Y=[]
[]	=	[X Y]	will always fail

Lists **Trees** Tutorial 2

Data structures: Tree

Tree below can be represented as either a+b*c or +(a,*(b,c))



Trees can be used to represent concept such as: sentence(noun(john),verb_phrase(verb(likes),noun(mary)))



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Family - rules

%Facts T2 father_of(tom, bill). mother_of(mary,jane).

```
%Rule T2
sibling(Person1,Person2):-
mother_of(Mother,Person1),
mother_of(Mother,Person2),
Person1\==Person2.
sibling(Person1,Person2):-
father_of(Father,Person1),
father_of(Father,Person2),
Person1\==Person2.
```

%Query T2 ?-sibling (anne, Person).



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Observations

- findall/3 predicate gathers all possible solutions for the query in argument 2 in a list.
- Procedural interpretation: to solve sibling/2, predicates mother_of/2 and father_of/2 must be evaluated first.
- The resolution mechanism (SLD resolution) builds an execution tree.

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Execution tree for sibling



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Section 3

Backtracking

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Obtaining multiple solutions in Prolog

- **Backtracking** in Prolog can be used to obtain multiple solutions for a goal (predicate fails vs user-induced failure).
- When Prolog is computing a goal, every choice it makes is stored as a **choice point**.
- If a given goal fails, it can be backtracked (choices undone) until the previous choice point is restored.
- From there, Prolog starts looking for a new alternative for solving the goal.

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Family - recursive

```
%Query T3
?-male_ancestor(P,anne,2).
```

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Observations

- Declarative interpretation of query: are **greg** and **anne** related through the male line?
- male_ancestor/3 can solve both the level of the relation and find a person who is related to someone else.
- Note that recursion must be made with new variables every time: **NewLevel**, **NewPerson**.

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Execution tree for male_ancestor



Section 4

Conclusion

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Conclusions

- Prolog allows for an almost direct match between a FOL algorithm and its execution syntax.
- Because Prolog is based on FOL, if the facts are true, then the program will produce true results.
- Prolog's declarative nature is particularly suited to build and consult multi-relational data.

There are many Prolog systems in the literature with different features such as Yap Prolog¹, SWI-Prolog², SICStus Prolog³, etc.

¹www.dcc.fc.up.pt/ vsc/Yap/documentation.html

²www.swi-prolog.org

³sicstus.sics.se

Thank you

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